

Safe-Error Attacks on SIKE and CSIDH

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- quantum computing \rightarrow impact on cryptography
- NIST PQC standardization process
- isogeny-based youngest field of post-quantum cryptography
- only little attention with respect to fault attack resilience
- we aim to fill this gap

Preliminaries

ECC vs Isogeny





Fundamental operation: $P \mapsto [n]P$

Nodes: curves over \mathbb{F}_{419} Edges: 3-, 5-, 7-isogenies Fundamental operation: $\varphi_A : E_0 \to E_a$

Isogeny graph mostly "stolen" from Chloe Martindale https://www.martindale.info/talks/QIT-Bristol.pdf

- key encapsulation
- over the quadratic extension field with $p = 2^n * 3^m \pm 1$
- alternate candidate at round 3 of NIST's (not a) competition
- four parameter sets: SIKEp434, SIKEp503, SIKEp610, SIKEp751.
- slow, but small key sizes

CSIDH (Commutative Supersingular Isogeny Diffie-Hellman)

- non-interactive key exchange protocol \rightarrow potential drop-in replacement for Diffie-Hellman
- over F_p with $p = 4 \cdot \ell_1 \cdots \ell_n 1$, where ℓ_1, \ldots, ℓ_n are small distinct odd primes
- private key = (e₁,...,e_n), where |e_i| = number of isogenies of degree ℓ_i (~ steps in the graph)
- not submitted to NIST's process designed after deadline
- post-quantum security under discussion
- even slower, but small key sizes

- adversary uses fault injections to perturb a specific location
- presence or absence of an error gives insight into the "codepath"
- memory (M) safe-error: the attacker modifies the memory
- computational (C) safe-error: computation attacked (skipping instructions)

Practical Experiments



Figure 1: ChipWhisperer cw1173

- ChipWhisperer-Lite ARM
- 32-bit STM32F303
- open source toolchain
- power analysis
- voltage and clock glitching

Setup

- attacker performs single fault injection per run
- · check if fault impacts shared secret
- · Bob's, Alice's public key, and shared secret precomputed
- · computation of shared secret attacked



- SIKEp434 Cortex-M4 implementation available at pqm4¹
- vulnerability remains the same across all available implementations
- 21,800 fault injections (100 injections for each bit)
- · injections during computation of the three-point ladder
- number of injections for full key recovery only depends on length
 of private key
- critical spots empirically determined with manageable effort

¹https://github.com/mupq/pqm4

```
1
          [...]
 2
         // main loop
 3
          for (i = 0; i < nbits; i++) {</pre>
 4
               bit = (m[i \rightarrow LOG2RADIX] \rightarrow (i \& (RADIX - 1))) \& 1;
 5
               swap = bit ^ prevbit;
 6
               prevbit = bit;
 7
               mask = 0 - (digit_t)swap;
8
9
               swap_points(R, R2, mask);
10
               xDBLADD(R0, R2, R-X, A24);
11
               fp2mul_mont(R2 \rightarrow X, R \rightarrow Z, R2 \rightarrow X);
12
13
          [...]
```

Listing 1: LADDER3PT@SIKE

- reduced the key space in CSIDH512 from $11^{74}~{\rm to}~3^2$
- attacker aims to distinguish a real from a dummy isogeny
- degree of attacked isogeny recovered with manageable effort
- 5,000 fault injections with high accuracy
- injections during the isogeny computation
- number of injections depends on many aspects

key	# of trials	faulty shared secret	accuracy
$S_1 = (-1, 1)$	2500	0.0%	100.0%
$S_2 = (0, 1)$	2500	92.4%	92.4%

Table 1: Results for CSIDH attacking the first isogeny

```
2/2
```

```
1
         [...]
2
        bool xISOG(proj *A, ..., int mask)
3
        {
 4
             proj Acopy = *A;
 5
             [...]
6
            // calculate new curve A
7
             [...]
8
             // CONSTANT TIME : swap back
9
             fp_cswap(&A->x, &Acopy.x, mask);
10
             fp_cswap(&A->z, &Acopy.z, mask);
11
             [...]
12
13
         [...]
```

Listing 2: xISOG@CSIDH



SIKE

- 5 fault injections at each bit leads to success rate above 99%
- full key recovery requires about 4 hours

CSIDH/MCR

- 1184 injections required for full key recovery
- full key recovery requires about 98 hours

- no big scandal
- securing cryptosystems against safe-error attacks is non-trivial
- dummy-free implementations of CSIDH, not vulnerable to the attacks
- more effort into the cryptanalysis of post-quantum candidates
- ChipWhisperer: perfectly adequate

Thank you for your attention!

Paper: https://eprint.iacr.org/2021/1132 Code: https://github.com/Safe-Error-Attacks-on-SIKE-and-CSIDH/SEAoSaC